

The Gabriel-Quillen functor for extriangulated categories

Yasuaki Ogawa

Nagoya University

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Notation

Categories and functors are always assumed to be additive.

- \mathcal{C} - category.
- $\text{Mod } \mathcal{C}$ - category of contravariant functors: $\mathcal{C} \rightarrow \text{Ab}$.

$$\begin{array}{ccc} \text{Mod } \mathcal{C} & \supseteq & \text{Lex } \mathcal{C} \\ \cup & & \cup \\ \text{mod } \mathcal{C} & \supseteq & \text{lex } \mathcal{C} \end{array}$$

- The Yoneda functor is denoted by $\mathbb{Y} : \mathcal{C} \hookrightarrow \text{Mod } \mathcal{C}$.

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$$\begin{array}{ccc} \text{Mod } \mathcal{C} & \supseteq & \text{Lex } \mathcal{C} \\ & & \text{left exact} \\ \cup & & \cup \\ \text{mod } \mathcal{C} & \supseteq & \text{lex } \mathcal{C} \\ \text{finitely presented} & & \end{array}$$

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- 2 Extriangulated category
- 3 The GQ functor for extriangulated categories
- 4 The general heart construction

The Gabriel-Quillen embedding theorem

Theorem

Let $(\mathcal{C}, \mathbb{E})$ be a small exact category.

- 1 There exists the following localization sequence

$$\text{Mod } \mathcal{C} \quad \begin{array}{c} \xrightarrow{Q} \\ \curvearrowleft R \end{array} \quad \text{Lex } \mathcal{C}$$

where R is the canonical inclusion.

- 2 The composed functor $E_{\mathcal{C}} : \mathcal{C} \hookrightarrow \text{Mod } \mathcal{C} \xrightarrow{Q} \text{Lex } \mathcal{C}$ is exact and fully faithful.
- 3 \mathcal{C} is extension-closed in $\text{Lex } \mathcal{C}$ and the exact structure \mathbb{E} is a class of all short sequences which belong to \mathcal{C} .

The Gabriel-Quillen embedding theorem

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The Gabriel-Quillen embedding theorem

Theorem

Let $(\mathcal{C}, \mathbb{E})$ be a small exact category.

- There exists the following localization sequence

$$\text{Ker } Q \begin{array}{c} \xrightarrow{\quad} \\ \curvearrowright \end{array} \text{Mod } \mathcal{C} \begin{array}{c} \xrightarrow{Q} \\ \curvearrowleft \\ R \end{array} \text{Lex } \mathcal{C}$$

where R is the canonical inclusion.

- The composed functor $E_{\mathcal{C}} : \mathcal{C} \hookrightarrow \text{Mod } \mathcal{C} \xrightarrow{Q} \text{Lex } \mathcal{C}$ is exact and fully faithful.
- \mathcal{C} is extension-closed in $\text{Lex } \mathcal{C}$ and the exact structure \mathbb{E} is a class of all short sequences which belong to \mathcal{C} .

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 & & \begin{array}{c} \uparrow \mathbb{Y} \\ \mathcal{C} \end{array} & & \\
 & & & \begin{array}{c} \curvearrowright \\ E_{\mathcal{C}} \end{array} &
 \end{array}$$

where R is the canonical inclusion.

- ② The composed functor $E_{\mathcal{C}} : \mathcal{C} \hookrightarrow \text{Mod } \mathcal{C} \xrightarrow{Q} \text{Lex } \mathcal{C}$ is exact and fully faithful.
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 & & \begin{array}{c} \uparrow \mathbb{Y} \\ \mathcal{C} \end{array} & & \begin{array}{c} \nearrow \\ E_{\mathcal{C}} \end{array}
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Closer look at the GQ theorem

Definition (Auslander)

Let $\delta : Z \rightarrow Y \rightarrow X$ be a conflation in \mathcal{C} . Then we have an exact sequence

$$0 \longrightarrow \mathcal{C}(-, Z) \longrightarrow \mathcal{C}(-, Y) \longrightarrow \mathcal{C}(-, X)$$

in $\text{Mod } \mathcal{C}$.

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defect of δ

in $\text{Mod } \mathcal{C}$.

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Definition (Auslander'66)

Define the following subcategories in $\text{Mod } \mathcal{C}$.

- $\text{def } \mathcal{C}$ - all defects in $\text{Mod } \mathcal{C}$.
- $\text{Def } \mathcal{C}$ - all filtered colimits of factors of defects in $\text{Mod } \mathcal{C}$.

Theorem

Let $(\mathcal{C}, \mathbb{E})$ be a small exact category. There exists the following localization sequence

$$\begin{array}{ccc}
 \xrightarrow{\quad} & \text{Mod } \mathcal{C} & \xrightarrow{Q} \\
 \curvearrowright & & \curvearrowleft \\
 & & R
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where R is the canonical inclusion.

The definition

The notion of extriangulated category is a simultaneous generalization of triangulated ones and exact ones.

Definition (Nakaoka-Palu'19)

The extriangulated category is a triple $(\mathcal{C}, \mathbb{E}, \mathfrak{s})$, where

- \mathcal{C} - an additive category;
- \mathbb{E} - a biadditive functor $\mathcal{C} \times \mathcal{C}^{\text{op}} \rightarrow \mathbf{Ab}$;
- \mathfrak{s} - an assignment from an element in $\mathbb{E}(X, Z)$ to a sequence $Z \rightarrow Y \rightarrow X$ in \mathcal{C} ,

with some suitable compatibility.

Some examples

- Extension-closed subcategory in a triangulated category has a natural extriangulated structure.
- Let $(\mathcal{U}, \mathcal{V})$ be a torsion pair in a triangulated category. Then \mathcal{U} has a natural extriangulated structure.

The GQ type localization sequence

Theorem (O)

Let \mathcal{C} be a small extriangulated category. Then there exists a localization sequence

$$\text{Def } \mathcal{C} \quad \begin{array}{c} \xrightarrow{\quad} \\ \curvearrowright \end{array} \quad \text{Mod } \mathcal{C} \quad \begin{array}{c} \xrightarrow{Q} \\ \curvearrowright \\ R \end{array} \quad \text{Lex } \mathcal{C}$$

where R is the canonical inclusion.

A finitely presented version

Lemma (Freyd'65)

The following are equivalent for an additive category \mathcal{C} :

- ① The category \mathcal{C} admits weak-kernels;
- ② The full subcategory $\text{mod } \mathcal{C}$ is an exact abelian subcategory in $\text{Mod } \mathcal{C}$.

Proposition (O)

Let \mathcal{C} be an extriangulated category with weak-kernels. Then $\text{def } \mathcal{C}$ is a Serre subcategory in $\text{mod } \mathcal{C}$.

Remark

The quotient functor $Q : \text{mod } \mathcal{C} \rightarrow \frac{\text{mod } \mathcal{C}}{\text{def } \mathcal{C}}$ does not necessarily have a right adjoint.

The GQ functor

This situation is depicted as follows:

$$\text{def } \mathcal{C} \longrightarrow \text{mod } \mathcal{C} \xrightarrow{Q} \frac{\text{mod } \mathcal{C}}{\text{def } \mathcal{C}}$$

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$$\text{def } \mathcal{C} \quad \xrightarrow{\quad} \quad \text{mod } \mathcal{C} \quad \xrightarrow[\circlearrowleft]{Q} \quad \frac{\text{mod } \mathcal{C}}{\text{def } \mathcal{C}}$$

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$$\begin{array}{ccc}
 \text{def } \mathcal{C} & \xrightarrow{\quad} & \text{mod } \mathcal{C} \\
 & \dashleftarrow \text{X} & \\
 & & \text{mod } \mathcal{C} \\
 & & \xrightarrow{Q} \\
 & & \dashleftarrow \text{X} \\
 & & \frac{\text{mod } \mathcal{C}}{\text{def } \mathcal{C}} \\
 & & \text{!!} \\
 & & \text{gq } \mathcal{C}
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 & \searrow \text{---} \times \text{---} & & \searrow \text{---} \times \text{---} & \\
 & & \mathbb{Y} & & \text{!!} \\
 & & \uparrow & & \\
 & & \mathcal{C} & \xrightarrow{Q \circ \mathbb{Y}} & \text{gq } \mathcal{C}
 \end{array}$$

We put $E_{\mathcal{C}} := Q \circ \mathbb{Y}$ and call it the *GQ functor* for \mathcal{C} .

A use of the GQ functor

Theorem (O)

The GQ functor $E_{\mathcal{C}} : \mathcal{C} \rightarrow \mathbf{gq}\mathcal{C}$ is:

- ① exact and fully faithful iff \mathcal{C} is an exact category;
- ② an equivalence iff \mathcal{C} is an abelian category.

Corollary (Auslander's defect formula)

Let \mathcal{C} be an abelian category. Then we have the localization sequence

$$\text{def } \mathcal{C} \quad \begin{array}{c} \xrightarrow{\quad} \\ \curvearrowright \end{array} \quad \text{mod } \mathcal{C} \quad \begin{array}{c} \xrightarrow{Q} \\ \curvearrowleft \\ R \end{array} \quad \mathbf{gq}\mathcal{C}$$

In particular, we have an equality $\mathcal{C} = \text{lex } \mathcal{C}$ in $\text{mod } \mathcal{C}$.

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In particular, we have an equality $\mathcal{C} = \text{lex } \mathcal{C}$ in $\text{mod } \mathcal{C}$.

The exact case

Theorem

Let \mathcal{C} be an idempotent complete exact category with weak-kernels. Assume that it has enough projectives.

- There exists the following localization sequence

$$\text{def } \mathcal{C} \quad \begin{array}{c} \xrightarrow{\quad} \\ \frown \end{array} \quad \text{mod } \mathcal{C} \quad \begin{array}{c} \xrightarrow{Q} \\ \smile \\ R \end{array} \quad \text{lex } \mathcal{C}$$

where R is the canonical inclusion.

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Torsion class

- \mathcal{T} - triangulated category.
- $(\mathcal{U}, \mathcal{V})$ - torsion pair.

Proposition

There exists the following localization sequence.

$$\text{def } \mathcal{U} \quad \begin{array}{c} \xrightarrow{\quad} \\ \curvearrowright \end{array} \quad \text{mod } \mathcal{U} \quad \begin{array}{c} \xrightarrow{Q} \\ \curvearrowright \\ R \end{array} \quad \text{lex } \mathcal{U}$$

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The heart of $(\mathcal{U}, \mathcal{V})$

- $(\mathcal{U}, \mathcal{V})$ - torsion pair in \mathcal{T} .
- Put $\mathcal{W} := \mathcal{U} \cap \mathcal{V}[-1]$.
- Put $\mathcal{T}^+ := \mathcal{U}[-1] * \mathcal{W}$, $\mathcal{T}^- := \mathcal{W} * \mathcal{V}$ and $\mathcal{H} := \mathcal{T}^+ \cap \mathcal{T}^-$.
- $\underline{\mathcal{H}} := \mathcal{H}/[\mathcal{W}]$ - the Nakaoka heart.

Theorem (Nakaoka)

The Nakaoka heart $\underline{\mathcal{H}}$ is abelian.

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Example

Let $(\mathcal{U}, \mathcal{V})$ be a torsion pair in \mathcal{T} .

- If $(\mathcal{U}, \mathcal{V}[-1])$ forms a t-structure, then $\underline{\mathcal{H}}$ is the usual heart of it.
- If \mathcal{U} is a (2-)cluster tilting subcategory, then $\underline{\mathcal{H}}$ is $\mathcal{T}/[\mathcal{C}]$ which has been shown to be abelian by Koenig-Zhu.

Theorem (O)

The Nakaoka heart $\underline{\mathcal{H}}$ is equivalent to $\text{lex}\mathcal{U}$.

Thank you for your attention!